

Definitions of Petri Net. Part I

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Summary. In the paper the classical definition of Petri net is described. The article also contains some theorems needed for proving equivalences of these definitions with other definitions of Petri net as relational algebras. See [3], [4] and other.

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The articles [5], [1], [6], and [2] provide the notation and terminology for this paper.

In this paper x, y, X, Y denote sets.

Let N be a net. The functor $\text{chaos}(N)$ yields a set and is defined by:

(Def. 2)¹ $\text{chaos}(N) = \text{Elements}(N) \cup \{\text{Elements}(N)\}$.

In the sequel M is a Petri net.

Let us consider X, Y . Let us assume that X misses Y . The functor $\text{PTempty}_f(X, Y)$ yielding a strict Petri net is defined by:

(Def. 4)² $\text{PTempty}_f(X, Y) = \langle X, Y, \emptyset \rangle$.

Let us consider X . The functor $\text{Tempy}_f(X)$ yields a strict Petri net and is defined by:

(Def. 5) $\text{Tempy}_f(X) = \text{PTempty}_f(X, \emptyset)$.

The functor $\text{Pempty}_f(X)$ yields a strict Petri net and is defined by:

(Def. 6) $\text{Pempty}_f(X) = \text{PTempty}_f(\emptyset, X)$.

Let us consider x . The functor $\text{Tsingle}_f(x)$ yields a strict Petri net and is defined as follows:

(Def. 7) $\text{Tsingle}_f(x) = \text{PTempty}_f(\emptyset, \{x\})$.

The functor $\text{Psingle}_f(x)$ yields a strict Petri net and is defined by:

(Def. 8) $\text{Psingle}_f(x) = \text{PTempty}_f(\{x\}, \emptyset)$.

The strict Petri net empty_f is defined as follows:

(Def. 9) $\text{empty}_f = \text{PTempty}_f(\emptyset, \emptyset)$.

One can prove the following propositions:

¹ The definition (Def. 1) has been removed.

² The definition (Def. 3) has been removed.

$\text{id}_{\text{Elements}(M)} = (\text{the flow relation of } M) \upharpoonright (\text{the places of } M)$ and $((\text{the flow relation of } M)^\smile \upharpoonright \text{the places of } M \cup \text{id}_{\text{the transitions of } M}) \setminus \text{id}_{\text{Elements}(M)} = (\text{the flow relation of } M)^\smile \upharpoonright \text{the places of } M$ and $((\text{the flow relation of } M) \upharpoonright (\text{the places of } M) \cup \text{id}_{\text{the transitions of } M}) \setminus \text{id}_{\text{Elements}(M)} = (\text{the flow relation of } M) \upharpoonright (\text{the places of } M)$ and $((\text{the flow relation of } M)^\smile \upharpoonright \text{the transitions of } M \cup \text{id}_{\text{the transitions of } M}) \setminus \text{id}_{\text{Elements}(M)} = (\text{the flow relation of } M)^\smile \upharpoonright \text{the transitions of } M$ and $((\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) \cup \text{id}_{\text{the transitions of } M}) \setminus \text{id}_{\text{Elements}(M)} = (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M)$.

- (20)(i) $((\text{The flow relation of } M) \upharpoonright (\text{the places of } M))^\smile = (\text{the flow relation of } M)^\smile \upharpoonright \text{the transitions of } M$, and
(ii) $((\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M))^\smile = (\text{the flow relation of } M)^\smile \upharpoonright \text{the places of } M$.
- (21)(i) $(\text{The flow relation of } M) \upharpoonright (\text{the places of } M) \cup (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) = \text{the flow relation of } M$,
(ii) $(\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) \cup (\text{the flow relation of } M) \upharpoonright (\text{the places of } M) = \text{the flow relation of } M$,
(iii) $((\text{the flow relation of } M) \upharpoonright (\text{the places of } M))^\smile \cup ((\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M))^\smile = (\text{the flow relation of } M)^\smile$, and
(iv) $((\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M))^\smile \cup ((\text{the flow relation of } M) \upharpoonright (\text{the places of } M))^\smile = (\text{the flow relation of } M)^\smile$.

Let us consider M . The functor $\text{enter}_f(M)$ yields a binary relation and is defined by:

(Def. 10) $\text{enter}_f(M) = (\text{the flow relation of } M)^\smile \upharpoonright \text{the transitions of } M \cup \text{id}_{\text{the places of } M}$.

The functor $\text{exit}_f(M)$ yields a binary relation and is defined as follows:

(Def. 11) $\text{exit}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) \cup \text{id}_{\text{the places of } M}$.

The following four propositions are true:

- (22) $\text{exit}_f(M) \subseteq [; \text{Elements}(M), \text{Elements}(M) ;]$ and $\text{enter}_f(M) \subseteq [; \text{Elements}(M), \text{Elements}(M) ;]$.
(23) $\text{dom exit}_f(M) \subseteq \text{Elements}(M)$ and $\text{rng exit}_f(M) \subseteq \text{Elements}(M)$ and $\text{dom enter}_f(M) \subseteq \text{Elements}(M)$ and $\text{rng enter}_f(M) \subseteq \text{Elements}(M)$.
(24) $\text{exit}_f(M) \cdot \text{exit}_f(M) = \text{exit}_f(M)$ and $\text{exit}_f(M) \cdot \text{enter}_f(M) = \text{exit}_f(M)$ and $\text{enter}_f(M) \cdot \text{enter}_f(M) = \text{enter}_f(M)$ and $\text{enter}_f(M) \cdot \text{exit}_f(M) = \text{enter}_f(M)$.
(25) $\text{exit}_f(M) \cdot (\text{exit}_f(M) \setminus \text{id}_{\text{Elements}(M)}) = \emptyset$ and $\text{enter}_f(M) \cdot (\text{enter}_f(M) \setminus \text{id}_{\text{Elements}(M)}) = \emptyset$.

Let us consider M . The functor $\text{prox}_f(M)$ yielding a binary relation is defined by the condition (Def. 12).

(Def. 12) $\text{prox}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the places of } M) \cup (\text{the flow relation of } M)^\smile \upharpoonright \text{the places of } M \cup \text{id}_{\text{the places of } M}$.

The functor $\text{flow}_f(M)$ yielding a binary relation is defined by:

(Def. 13) $\text{flow}_f(M) = (\text{the flow relation of } M) \cup \text{id}_{\text{Elements}(M)}$.

We now state the proposition

- (26) $\text{prox}_f(M) \cdot \text{prox}_f(M) = \text{prox}_f(M)$ and $(\text{prox}_f(M) \setminus \text{id}_{\text{Elements}(M)}) \cdot \text{prox}_f(M) = \emptyset$ and $\text{prox}_f(M) \cup (\text{prox}_f(M))^\smile \cup \text{id}_{\text{Elements}(M)} = \text{flow}_f(M) \cup (\text{flow}_f(M))^\smile$.

Let us consider M . The functor $\text{places}_f(M)$ yielding a set is defined as follows:

(Def. 14) $\text{places}_f(M) = \text{the places of } M$.

The functor $\text{transitions}_f(M)$ yielding a set is defined as follows:

(Def. 15) $\text{transitions}_f(M) =$ the transitions of M .

The functor $\text{pre}_f(M)$ yielding a binary relation is defined by:

(Def. 16) $\text{pre}_f(M) =$ (the flow relation of M) \upharpoonright (the transitions of M).

The functor $\text{post}_f(M)$ yielding a binary relation is defined by:

(Def. 17) $\text{post}_f(M) =$ (the flow relation of M) \smile \upharpoonright the transitions of M .

Next we state three propositions:

(27) $\text{pre}_f(M) \subseteq [:\text{transitions}_f(M), \text{places}_f(M):]$ and $\text{post}_f(M) \subseteq [:\text{transitions}_f(M), \text{places}_f(M):]$.

(28) $\text{places}_f(M)$ misses $\text{transitions}_f(M)$.

(29) $\text{prox}_f(M) \subseteq [:\text{Elements}(M), \text{Elements}(M):]$ and $\text{flow}_f(M) \subseteq [:\text{Elements}(M), \text{Elements}(M):]$.

Let us consider M . The functor $\text{entrance}_f(M)$ yielding a binary relation is defined as follows:

(Def. 18) $\text{entrance}_f(M) =$ (the flow relation of M) \smile \upharpoonright the places of $M \cup \text{id}_{\text{the transitions of } M}$.

The functor $\text{escape}_f(M)$ yielding a binary relation is defined as follows:

(Def. 19) $\text{escape}_f(M) =$ (the flow relation of M) \upharpoonright (the places of M) $\cup \text{id}_{\text{the transitions of } M}$.

Next we state four propositions:

(30) $\text{escape}_f(M) \subseteq [:\text{Elements}(M), \text{Elements}(M):]$ and $\text{entrance}_f(M) \subseteq [:\text{Elements}(M), \text{Elements}(M):]$.

(31) $\text{dom escape}_f(M) \subseteq \text{Elements}(M)$ and $\text{rng escape}_f(M) \subseteq \text{Elements}(M)$ and $\text{dom entrance}_f(M) \subseteq \text{Elements}(M)$ and $\text{rng entrance}_f(M) \subseteq \text{Elements}(M)$.

(32) $\text{escape}_f(M) \cdot \text{escape}_f(M) = \text{escape}_f(M)$ and $\text{escape}_f(M) \cdot \text{entrance}_f(M) = \text{escape}_f(M)$ and $\text{entrance}_f(M) \cdot \text{entrance}_f(M) = \text{entrance}_f(M)$ and $\text{entrance}_f(M) \cdot \text{escape}_f(M) = \text{entrance}_f(M)$.

(33) $\text{escape}_f(M) \cdot (\text{escape}_f(M) \setminus \text{id}_{\text{Elements}(M)}) = \emptyset$ and $\text{entrance}_f(M) \cdot (\text{entrance}_f(M) \setminus \text{id}_{\text{Elements}(M)}) = \emptyset$.

Let us consider M . The functor $\text{adjac}_f(M)$ yielding a binary relation is defined by the condition (Def. 20).

(Def. 20) $\text{adjac}_f(M) =$ (the flow relation of M) \upharpoonright (the transitions of M) \cup (the flow relation of M) \smile \upharpoonright the transitions of $M \cup \text{id}_{\text{the transitions of } M}$.

We introduce $\text{circulation}_f(M)$ as a synonym of $\text{flow}_f(M)$.

The following proposition is true

(34) $\text{adjac}_f(M) \cdot \text{adjac}_f(M) = \text{adjac}_f(M)$ and $(\text{adjac}_f(M) \setminus \text{id}_{\text{Elements}(M)}) \cdot \text{adjac}_f(M) = \emptyset$ and $\text{adjac}_f(M) \cup (\text{adjac}_f(M))^\smile \cup \text{id}_{\text{Elements}(M)} = \text{circulation}_f(M) \cup (\text{circulation}_f(M))^\smile$.

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