

On an Algorithmic Algebra over Simple-Named Complex-Valued Nominative Data

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Summary. This paper continues formalization in the Mizar system [2, 1] of basic notions of the composition-nominative approach to program semantics [14] which was started in [8, 12, 10].

The composition-nominative approach studies mathematical models of computer programs and data on various levels of abstraction and generality and provides tools for reasoning about their properties. In particular, data in computer systems are modeled as nominative data [15]. Besides formalization of semantics of programs, certain elements of the composition-nominative approach were applied to abstract systems in a mathematical systems theory [4, 6, 7, 5, 3].

In the paper we give a formal definition of the notions of a binominative function over given sets of names and values (i.e. a partial function which maps simple-named complex-valued nominative data to such data) and a nominative predicate (a partial predicate on simple-named complex-valued nominative data). The sets of such binominative functions and nominative predicates form the carrier of the generalized Glushkov algorithmic algebra for simple-named complex-valued nominative data [15]. This algebra can be used to formalize algorithms which operate on various data structures (such as multidimensional arrays, lists, etc.) and reason about their properties.

In particular, we formalize the operations of this algebra which require a specification of a data domain and which include the existential quantifier, the assignment composition, the composition of superposition into a predicate, the composition of superposition into a binominative function, the name checking

predicate. The details on formalization of nominative data and the operations of the algorithmic algebra over them are described in [11, 13, 9].

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1. PRELIMINARIES

From now on a, b, c, v, v_1, x, y denote objects, V, A denote sets, and d denotes a nominative data with simple names from V and complex values from A .

Now we state the proposition:

- (1) $\{a, b, c\} \subseteq A$ if and only if $a, b, c \in A$.

Let a, b, c, d, e, f be objects. One can verify that $\{\langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle\}$ is relation-like.

Let us consider objects a, b, c, d, e, f . Now we state the propositions:

- (2) $\text{dom}\{\langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle\} = \{a, c, e\}$.
 (3) $\text{rng}\{\langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle\} = \{b, d, f\}$.

Let us consider V . Note that there exists a finite sequence which is one-to-one and V -valued.

- (4) $\text{dom}\langle a, b, c \rangle = \{1, 2, 3\}$.

Let us consider V and A . Let us note that $\text{ND}_{\text{SS}}(V, A)$ is non empty and has not non empty elements and $\text{ND}_{\text{SC}}(V, A)$ is non empty and has not non empty elements.

Now we state the propositions:

- (5) If $v \in V$, then $\{\langle v, d \rangle\}$ is a non-atomic nominative data of V and A .
 (6) Let us consider a finite function D . Suppose $\text{dom } D \subseteq V$ and $\text{rng } D \subseteq \text{ND}_{\text{SC}}(V, A)$. Then D is a non-atomic nominative data of V and A .

PROOF: Define $\mathcal{P}[\text{set}] \equiv \mathcal{S}_1$ is a non-atomic nominative data of V and A . For every sets x, B such that $x \in D$ and $B \subseteq D$ and $\mathcal{P}[B]$ holds $\mathcal{P}[B \cup \{x\}]$. $\mathcal{P}[D]$. \square

- (7) Let us consider nominative data d_1, d_2 with simple names from V and complex values from A . Then $d_2 \subseteq d_1 \nabla_a d_2$.
 (8) Every non-atomic nominative data of V and A is a nominative data with simple names from V and complex values from A .

- (9) Let us consider non-atomic nominative data d_1, d_2 of V and A . Then $d_1 \nabla_a d_2$ is a non-atomic nominative data of V and A . The theorem is a consequence of (8).

Let us consider V and A . Let d_1, d_2 be non-atomic nominative data of V and A . Let us observe that $d_1 \nabla_a d_2$ is function-like and relation-like.

Let us consider v . One can verify that $d_1 \nabla_a^v d_2$ is function-like and relation-like.

Let d_1 be a non-atomic nominative data of V and A and d_2 be a nominative data with simple names from V and complex values from A . Let us observe that $d_1 \nabla_a^v d_2$ is function-like and relation-like.

Now we state the propositions:

- (10) Suppose $v \in V$. Let us consider nominative data d_1, d_2 with simple names from V and complex values from A , and a function L . If $L = d_1 \nabla_a^v d_2$, then $L(v) = d_2$. The theorem is a consequence of (8).
- (11) Suppose $v \in V$ and $v \neq v_1$. Let us consider a non-atomic nominative data d_1 of V and A , a nominative data d_2 with simple names from V and complex values from A , and a function L . Suppose $L = d_1 \nabla_a^v d_2$ and $v_1 \in \text{dom } d_1$ and $d_1 \notin A$ and $\Rightarrow v(d_2) \notin A$. Then $L(v_1) = d_1(v_1)$. The theorem is a consequence of (8).

Let us consider a non-atomic nominative data d_1 of V and A and a nominative data d_2 with simple names from V and complex values from A . Now we state the propositions:

- (12) Suppose $v \in V$ and $v \notin \text{dom } d_1$ and $d_1 \notin A$ and $\Rightarrow v(d_2) \notin A$. Then $\text{dom}(d_1 \nabla_a^v d_2) = \{v\} \cup \text{dom } d_1$.
- (13) If $v \in V$ and $v \in \text{dom } d_1$ and $d_1 \notin A$ and $\Rightarrow v(d_2) \notin A$, then $\text{dom}(d_1 \nabla_a^v d_2) = \text{dom } d_1$.
- (14) If $v \in V$ and $d_1 \notin A$ and $\Rightarrow v(d_2) \notin A$, then $\text{dom}(d_1 \nabla_a^v d_2) = \{v\} \cup \text{dom } d_1$.
The theorem is a consequence of (13) and (12).

Let us consider V and A .

A partial predicate over simple-named complex-valued nominative data of V and A is a partial predicate of $\text{ND}_{\text{SC}}(V, A)$. In the sequel p, q, r denote partial predicates over simple-named complex-valued nominative data of V and A .

Now we state the propositions:

- (15) $\text{dom}(p \vee q) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d \in \text{dom } p \text{ and } p(d) = \text{true} \text{ or } d \in \text{dom } q \text{ and } q(d) = \text{true} \text{ or } d \in \text{dom } p \text{ and } p(d) = \text{false} \text{ and } d \in \text{dom } q \text{ and } q(d) = \text{false}\}$.
- (16) $\text{dom}(p \wedge q) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d \in \text{dom } p \text{ and } p(d) = \text{false} \text{ or } d \in$

$\text{dom } q$ and $q(d) = \text{false}$ or $d \in \text{dom } p$ and $p(d) = \text{true}$ and $d \in \text{dom } q$ and $q(d) = \text{true}$ }.

- (17) $\text{dom}(p \Rightarrow q) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d \in \text{dom } p \text{ and } p(d) = \text{false} \text{ or } d \in \text{dom } q \text{ and } q(d) = \text{true} \text{ or } d \in \text{dom } p \text{ and } p(d) = \text{true} \text{ and } d \in \text{dom } q \text{ and } q(d) = \text{false}\}$.

Let us consider V , A , and v . The functor $\exists_v^{V,A}$ yielding a function from $\text{Pr}(\text{ND}_{\text{SC}}(V, A))$ into $\text{Pr}(\text{ND}_{\text{SC}}(V, A))$ is defined by

- (Def. 1) for every partial predicate over simple-named complex-valued nominative data p of V and A , $\text{dom}(it(p)) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : \text{ there exists a nominative data } d_1 \text{ with simple names from } V \text{ and complex values from } A \text{ such that } d\nabla_a^v d_1 \in \text{dom } p \text{ and } p(d\nabla_a^v d_1) = \text{true} \text{ or for every nominative data } d_1 \text{ with simple names from } V \text{ and complex values from } A, d\nabla_a^v d_1 \in \text{dom } p \text{ and } p(d\nabla_a^v d_1) = \text{false}\}$ and for every nominative data d with simple names from V and complex values from A , if there exists a nominative data d_1 with simple names from V and complex values from A such that $d\nabla_a^v d_1 \in \text{dom } p$ and $p(d\nabla_a^v d_1) = \text{true}$, then $it(p)(d) = \text{true}$ and if for every nominative data d_1 with simple names from V and complex values from A , $d\nabla_a^v d_1 \in \text{dom } p$ and $p(d\nabla_a^v d_1) = \text{false}$, then $it(p)(d) = \text{false}$.

Let us consider p . The functor $\exists_v p$ yielding a partial predicate over simple-named complex-valued nominative data of V and A is defined by the term

- (Def. 2) $(\exists_v^{V,A})(p)$.

Now we state the propositions:

- (18) Suppose $x \in \text{dom}(\exists_v p)$. Then

- (i) there exists a nominative data d_1 with simple names from V and complex values from A such that $x\nabla_a^v d_1 \in \text{dom } p$ and $p(x\nabla_a^v d_1) = \text{true}$, or
- (ii) for every nominative data d_1 with simple names from V and complex values from A , $x\nabla_a^v d_1 \in \text{dom } p$ and $p(x\nabla_a^v d_1) = \text{false}$.

- (19) $\exists_v \perp_{\text{PP}}(\text{ND}_{\text{SC}}(V, A)) = \perp_{\text{PP}}(\text{ND}_{\text{SC}}(V, A))$. The theorem is a consequence of (18).

- (20) DISTRIBUTIVITY LAW:

$$\exists_v(p \vee q) = \exists_v p \vee \exists_v q.$$

2. ON AN ALGORITHMIC ALGEBRA OVER SIMPLE-NAMED COMPLEX-VALUED NOMINATIVE DATA

From now on n denotes a natural number and X denotes a function.

Let F be a function yielding function and d be an object. We say that d is in $\text{doms } F$ if and only if

(Def. 3) for every object x such that $x \in \text{dom } F$ holds $d \in \text{dom}(F(x))$.

Let g be a function yielding function and X be a function. The functor $\text{NDdataSeq}(g, X, d)$ yielding a function is defined by

(Def. 4) $\text{dom } it = \text{dom } X$ and for every x such that $x \in \text{dom } X$ holds $it(x) = \langle X(x), g(x)(d) \rangle$.

Let X be a finite function. Let us note that $\text{NDdataSeq}(g, X, d)$ is finite.

Let X be a finite sequence. One can check that $\text{NDdataSeq}(g, X, d)$ is finite sequence-like.

Let X be a function. The functor $\text{NDentry}(g, X, d)$ yielding a set is defined by the term

(Def. 5) $\text{rng } \text{NDdataSeq}(g, X, d)$.

Now we state the propositions:

(21) Let us consider a function f , and objects a, d . Then $\text{NDentry}(\langle f \rangle, \langle a \rangle, d) = \{ \langle a, f(d) \rangle \}$.

(22) Let us consider functions f, g , and objects a, b, d . Then $\text{NDentry}(\langle f, g \rangle, \langle a, b \rangle, d) = \{ \langle a, f(d) \rangle, \langle b, g(d) \rangle \}$.

(23) Let us consider functions f, g, h , and objects a, b, c, d . Then $\text{NDentry}(\langle f, g, h \rangle, \langle a, b, c \rangle, d) = \{ \langle a, f(d) \rangle, \langle b, g(d) \rangle, \langle c, h(d) \rangle \}$. The theorem is a consequence of (4).

Let g be a function yielding function, X be a function, and d be an object. Let us note that $\text{NDentry}(g, X, d)$ is relation-like.

Let X be a one-to-one function. One can verify that $\text{NDentry}(g, X, d)$ is function-like.

Let X be a finite function. Observe that $\text{NDentry}(g, X, d)$ is finite.

Now we state the proposition:

(24) Let us consider a function yielding function g , a function X , and an object d . Then $\text{dom}(\text{NDentry}(g, X, d)) = \text{rng } X$.

Let us consider V and A .

A binominative function over simple-named complex-valued nominative data of V and A is a partial function from $\text{ND}_{\text{SC}}(V, A)$ to $\text{ND}_{\text{SC}}(V, A)$. From now on f, g, h denote binominative functions over simple-named complex-valued nominative data of V and A .

Now we state the propositions:

(25) $\text{rng NDdataSeq}(\langle f \rangle, \langle v \rangle, d) = v \mapsto f(d)$.

(26) If $a \in V$ and $d \in \text{dom } f$, then $\text{NDentry}(\langle f \rangle, \langle a \rangle, d) = \Rightarrow a(f(d))$. The theorem is a consequence of (25).

(27) If $a \in V$ and $d \in \text{dom } f$, then $\text{NDentry}(\langle f \rangle, \langle a \rangle, d)$ is a non-atomic nominative data of V and A . The theorem is a consequence of (26).

(28) Suppose $\{a, b\} \subseteq V$ and $a \neq b$ and $d \in \text{dom } f$ and $d \in \text{dom } g$. Then $\text{NDentry}(\langle f, g \rangle, \langle a, b \rangle, d)$ is a non-atomic nominative data of V and A . The theorem is a consequence of (22) and (6).

(29) Suppose $\{a, b, c\} \subseteq V$ and a, b, c are mutually different and $d \in \text{dom } f$ and $d \in \text{dom } g$ and $d \in \text{dom } h$. Then $\text{NDentry}(\langle f, g, h \rangle, \langle a, b, c \rangle, d)$ is a non-atomic nominative data of V and A . The theorem is a consequence of (23), (2), (3), (1), and (6).

Let us consider V and A . Let f be a finite sequence. We say that f is (V, A) -FPrg-yielding if and only if

(Def. 6) for every n such that $1 \leq n \leq \text{len } f$ holds $f(n)$ is a binominative function over simple-named complex-valued nominative data of V and A .

Let us consider f . Let us note that $\langle f \rangle$ is (V, A) -FPrg-yielding.

Let us consider g . Note that $\langle f, g \rangle$ is (V, A) -FPrg-yielding.

Let us consider h . Let us observe that $\langle f, g, h \rangle$ is (V, A) -FPrg-yielding.

Let us consider n . One can verify that there exists a finite sequence which is (V, A) -FPrg-yielding and n -element.

Let us consider x . Let g be a (V, A) -FPrg-yielding finite sequence. One can verify that $g(x)$ is function-like and relation-like and every finite sequence which is (V, A) -FPrg-yielding is also function yielding.

Now we state the propositions:

(30) Let us consider a (V, A) -FPrg-yielding finite sequence g , and a one-to-one finite sequence X . Suppose $\text{dom } g = \text{dom } X$ and d is in $\text{doms } g$. Then $\text{rng NDentry}(g, X, d) \subseteq \text{ND}_{\text{SC}}(V, A)$.

(31) Let us consider a (V, A) -FPrg-yielding finite sequence g , and a one-to-one, V -valued finite sequence X . Suppose $\text{dom } g = \text{dom } X$ and d is in $\text{doms } g$. Then $\text{NDentry}(g, X, d)$ is a non-atomic nominative data of V and A . The theorem is a consequence of (24), (30), and (6).

Let us consider V , A , and v . The functor $\text{Asg}^{V, A, v}$ yielding a function from $\text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ into $\text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ is defined by

(Def. 7) for every binominative function over simple-named complex-valued nominative data f of V and A , $\text{dom}(it(f)) = \text{dom } f$ and for every nominative

data d with simple names from V and complex values from A such that $d \in \text{dom}(it(f))$ holds $it(f)(d) = d\nabla_a^v f(d)$.

Let us consider V , A , v , and f . The functor $\text{Asg}^v(f)$ yielding a binominative function over simple-named complex-valued nominative data of V and A is defined by the term

(Def. 8) $\text{Asg}^{V,A,v}(f)$.

Let d be a non-atomic nominative data of V and A . One can check that $\text{Asg}^v(f)(d)$ is function-like and relation-like.

Now we state the proposition:

(32) Let us consider a non-atomic nominative data d of V and A . Suppose $v \in V$ and $d \notin A$ and $\Rightarrow v(f(d)) \notin A$ and $d \in \text{dom } f$. Then $\text{dom}((\text{Asg}^v(f))(d)) = \text{dom } d \cup \{v\}$. The theorem is a consequence of (14).

Let us consider V and A . Let g be a (V,A) -FPrg-yielding finite sequence. Assume $\prod g \neq \emptyset$. Let X be a function. The functor $\text{S}_P(g, X)$ yielding a function from $\text{Pr}(\text{ND}_{\text{SC}}(V, A)) \times \prod g$ into $\text{Pr}(\text{ND}_{\text{SC}}(V, A))$ is defined by

(Def. 9) for every partial predicate over simple-named complex-valued nominative data p of V and A and for every element x of $\prod g$, $\text{dom } it(p, x) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d\nabla_a(\text{NDentry}(g, X, d)) \in \text{dom } p \text{ and } d \text{ is in doms } g\}$ and for every nominative data d with simple names from V and complex values from A such that d is in doms g holds $it(p, x)(d) \cong p(d\nabla_a(\text{NDentry}(g, X, d)))$.

Let us consider V , A , and p . Let g be a (V,A) -FPrg-yielding finite sequence. Assume $\prod g \neq \emptyset$. Let X be a function and x be an element of $\prod g$. The functor $\text{S}_P(p, x, X)$ yielding a partial predicate over simple-named complex-valued nominative data of V and A is defined by the term

(Def. 10) $\text{S}_P(g, X)(p, x)$.

Now we state the proposition:

(33) Let us consider a (V,A) -FPrg-yielding finite sequence g . Suppose $\prod g \neq \emptyset$. Let us consider an element x of $\prod g$. Suppose $d \in \text{dom}(\text{S}_P(p, x, X))$. Then

- (i) d is in doms g , and
- (ii) $\text{S}_P(p, x, X)(d) = p(d\nabla_a(\text{NDentry}(g, X, d)))$.

Let us consider V , A , and v . The functor $\text{S}_P^{V,A,v}$ yielding a function from $\text{Pr}(\text{ND}_{\text{SC}}(V, A)) \times \text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ into $\text{Pr}(\text{ND}_{\text{SC}}(V, A))$ is defined by

(Def. 11) for every partial predicate over simple-named complex-valued nominative data p of V and A and for every binominative function over simple-named complex-valued nominative data f of V and A , $\text{dom } it(p, f) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex}$

values from $A : d\nabla_a^v f(d) \in \text{dom } p$ and $d \in \text{dom } f$ and for every nominative data d with simple names from V and complex values from A such that $d \in \text{dom } f$ holds $it(p, f)(d) \cong p(d\nabla_a^v f(d))$.

Let us consider V , A , v , p , and f . The functor $S_P(p, f, v)$ yielding a partial predicate over simple-named complex-valued nominative data of V and A is defined by the term

(Def. 12) $(S_P^{V,A,v})(p, f)$.

Now we state the propositions:

(34) If $d \in \text{dom}(S_P(p, f, v))$, then $S_P(p, f, v)(d) = p(d\nabla_a^v f(d))$ and $d \in \text{dom } f$.

(35) Let us consider an element x of $\prod \langle f \rangle$. Suppose $v \in V$ and $\prod \langle f \rangle \neq \emptyset$. Then $S_P(p, f, v) = S_P(p, x, \langle v \rangle)$. The theorem is a consequence of (26), (33), and (34).

Let us consider V and A . Let g be a (V,A) -FPrg-yielding finite sequence. Assume $\prod g \neq \emptyset$. Let X be a function. The functor $S_F(g, X)$ yielding a function from $\text{FPrg}(\text{ND}_{\text{SC}}(V, A)) \times \prod g$ into $\text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ is defined by

(Def. 13) for every binominative function over simple-named complex-valued nominative data f of V and A and for every element x of $\prod g$, $\text{dom } it(f, x) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d\nabla_a(\text{NDentry}(g, X, d)) \in \text{dom } f \text{ and } d \text{ is in doms } g\}$ and for every nominative data d with simple names from V and complex values from A such that d is in doms g holds $it(f, x)(d) \cong f(d\nabla_a(\text{NDentry}(g, X, d)))$.

Let us consider V , A , and f . Let g be a (V,A) -FPrg-yielding finite sequence. Assume $\prod g \neq \emptyset$. Let X be a function and x be an element of $\prod g$. The functor $S_F(f, x, X)$ yielding a binominative function over simple-named complex-valued nominative data of V and A is defined by the term

(Def. 14) $S_F(g, X)(f, x)$.

Now we state the proposition:

(36) Let us consider a (V,A) -FPrg-yielding finite sequence g . Suppose $\prod g \neq \emptyset$. Let us consider an element x of $\prod g$. Suppose $d \in \text{dom}(S_F(f, x, X))$. Then

(i) d is in doms g , and

(ii) $S_F(f, x, X)(d) = f(d\nabla_a(\text{NDentry}(g, X, d)))$.

Let us consider V , A , and v . The functor $S_F^{V,A,v}$ yielding a function from $\text{FPrg}(\text{ND}_{\text{SC}}(V, A)) \times \text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ into $\text{FPrg}(\text{ND}_{\text{SC}}(V, A))$ is defined by

(Def. 15) for every binominative functions over simple-named complex-valued nominative data f, g of V and A , $\text{dom } it(f, g) = \{d, \text{ where } d \text{ is a nominative data with simple names from } V \text{ and complex values from } A : d\nabla_a^v g(d) \in$

$\text{dom } f$ and $d \in \text{dom } g$ and for every nominative data d with simple names from V and complex values from A such that $d \in \text{dom } g$ holds $it(f, g)(d) \cong f(d\nabla_a^v g(d))$.

Let us consider V , A , v , f , and g . The functor $S_F(f, g, v)$ yielding a binominative function over simple-named complex-valued nominative data of V and A is defined by the term

(Def. 16) $(S_F^{V,A,v})(f, g)$.

Now we state the propositions:

(37) If $d \in \text{dom}(S_F(f, g, v))$, then $S_F(f, g, v)(d) = f(d\nabla_a^v g(d))$ and $d \in \text{dom } g$.

(38) Let us consider an element x of $\Pi\langle g \rangle$. Suppose $v \in V$ and $\Pi\langle g \rangle \neq \emptyset$. Then $S_F(f, g, v) = S_F(f, x, \langle v \rangle)$. The theorem is a consequence of (26), (36), and (37).

Let us consider V , A , and v . The functor $v!^{V,A}$ yielding a partial predicate over simple-named complex-valued nominative data of V and A is defined by

(Def. 17) $\text{dom } it = \text{ND}_{\text{SC}}(V, A) \setminus A$ and for every non-atomic nominative data d of V and A such that $d \in \text{dom } it$ holds if $v \Rightarrow_a d \in \text{dom } it$, then $it(d) = \text{true}$ and if $v \Rightarrow_a d \notin \text{dom } it$, then $it(d) = \text{false}$.

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